

Finite posets as algebraic expressions in duoidal categories

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Octoberfest 2023



Duoidal categories

A *duoidal* category \mathcal{C} consists of

- two monoidal structures (I, \otimes) and (J, \triangleleft) ;
- lax monoidal structures on the functors $\triangleleft: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ and $J: 1 \rightarrow \mathcal{C}$ with respect to (I, \otimes) (+ properties)
- equivalently, natural interchange morphisms (+ properties)

$$(a \triangleleft b) \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft (b \otimes d)$$

$$I \rightarrow I \triangleleft I \qquad J \otimes J \rightarrow J \qquad I \rightarrow J$$

A *physical* duoidal category moreover has

- \otimes is symmetric (compatibly with interchange)
- $I \cong J$ (compatibly with interchange)

$$\left[(P \triangleleft Q) \otimes (R \triangleleft S) \rightarrow (P \otimes R) \triangleleft (Q \otimes S), \quad \otimes \text{ symmetric, } I \cong J \right]$$

The category of posets has a physical duoidal structure:

- The unit is the empty poset \emptyset
- \otimes is disjoint union, \triangleleft is *join*

$$P \otimes Q = \left(P \quad Q \right) \quad P \triangleleft Q = \begin{pmatrix} Q \\ \uparrow \\ P \end{pmatrix}$$

where $P \triangleleft Q$ has elements $P \sqcup Q$ and $p < q$ for $p \in P, q \in Q$

$$\bullet \quad \begin{pmatrix} Q & S \\ \uparrow & \uparrow \\ P & R \end{pmatrix} \rightsquigarrow \begin{pmatrix} Q & S \\ \uparrow & \uparrow \\ P & R \end{pmatrix}$$

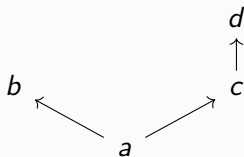
(Note: The diagram on the right shows two crossing arrows between the bottom and top levels, representing the symmetric property of the tensor product.)

Expressions

$$\left[(a \triangleleft b) \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft (b \otimes d), \quad \otimes \text{ symmetric, } I \cong J \right]$$

Physical duoidal categories have many compound operations and natural maps between them:

$$a \triangleleft (b \otimes (c \triangleleft d))$$



$$a \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft d \quad \left(\begin{array}{c} d \\ \uparrow \\ a \end{array} \right) \rightsquigarrow \left(\begin{array}{c} d \\ \uparrow \\ a \end{array} \right)$$

(S.–Spivak) These operations and natural maps correspond to *expressible* posets and bijective monotone functions.

A poset is expressible if it is either empty or can be constructed from singleton posets by finitely many disjoint unions and joins.

Expressible posets

$$\left[(a \triangleleft b) \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft (b \otimes d), \quad \otimes \text{ symmetric}, \quad I \cong J \right]$$

A poset is expressible if it is either empty or can be constructed from singleton posets by finitely many disjoint unions and joins.

Not all finite posets are expressible:

$$Z = \left(\begin{array}{ccc} & b & d \\ & \uparrow & \uparrow \\ & a & c \\ & & \swarrow \end{array} \right)$$

(S.–Spivak) A poset is expressible if and only if it is finite and admits no full embeddings of Z .

A categorical operad for physical duoidal categories

$$\left[(a \triangleleft b) \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft (b \otimes d), \quad \otimes \text{ symmetric, } I \cong J \right]$$

There is a categorical symmetric operad **Expr** where **Expr**_{*n*} is the category of expressible posets with *n* elements and bijective monotone functions. The unit is the singleton poset and operadic composition is given by “substitution”

$$\left((a \otimes c) \triangleleft (b \otimes d) \right); P, Q, R, S \quad \mapsto \quad \left(\begin{array}{ccc} Q & & S \\ \uparrow & \swarrow & \nearrow \\ P & & R \end{array} \right)$$

An **Expr**-algebra is a category \mathcal{C} with coherent functors

$$\mathbf{Expr}_n \times \mathcal{C}^n \rightarrow \mathcal{C}.$$

(S.–Spivak) **Expr**-algebras are precisely physical duoidal categories.

Nonnegative reals and parallel programs

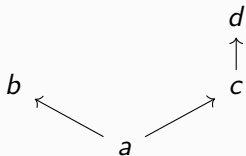
$$\left[(a \triangleleft b) \otimes (c \triangleleft d) \rightarrow (a \otimes c) \triangleleft (b \otimes d), \quad \otimes \text{ symmetric, } I \cong J \right]$$

The posetal category $\mathbb{R}_{\geq 0}$ of nonnegative real numbers has a physical duoidal structure with unit 0, \otimes given by maximum and \triangleleft given by addition.

$$(a + b) \max (c + d) \leq (a \max c) + (b \max d)$$

max and $+$ correspond to how the runtimes of two programs behave when they are run in parallel and series, respectively.

Given a finite set of programs and a poset of dependencies between them, the corresponding operation in $\mathbb{R}_{\geq 0}$ computes its optimal runtime.



- Brandon T. Shapiro and David I. Spivak, “Duoidal Structures for Compositional Dependence” arXiv:2210.01962

Thanks!